MTH 995-003: Intro to CS and Big Data

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## 1 The R.I.P. and Basis pursuit

**Definition 1.** A matrix  $A \in \mathbb{R}^{m \times N}$  satisfies the restricted isometry property of order  $k \in [N]$  for  $\epsilon \in (0,1)$  if

$$(1-\epsilon) \| \vec{x} \|_2^2 \le \|A\vec{x}\|_2^2 \le (1+\epsilon) \|\vec{x}\|_2^2 \tag{1}$$

for all vectors  $\vec{x} \in \mathbb{R}^N$  with  $\|\vec{x}\|_0 \leq k$  (at most k nonzero entries)

**Theorem 1.** let  $p, \epsilon \in (0,1)$  and  $A \in \mathbb{R}^{m \times N}$  with *i.i.d mean 0, variance 1, subgaussian entries* (with parameter c), and choose

$$m \geq \frac{32c(16c+1)}{(\sqrt{1+\epsilon}-1)^2} \cdot k \cdot \ln\left(\frac{eN\left(1+\frac{16}{\sqrt{1+\epsilon}-1}\right)}{k \cdot \sqrt[k]{(1-p)/2}}\right).$$

Then  $\frac{1}{\sqrt{m}}A$  will have RIP of order k for  $\epsilon$  with probability at least p.

*Proof:* Do as Homework problem 1.

**Definition 2.** Let  $\vec{x} \in \mathbb{R}^N$ . We let the support of  $\vec{x}$  be

 $supp(\vec{x}) = The set of nonzero coordinates in \vec{x}.$ 

Example 1. If  $\vec{x} = \begin{pmatrix} 100 \\ 0 \\ 100 \\ 0 \end{pmatrix}$ , then  $supp(\vec{x}) = \{1, 3\}$ .

**Definition 3.** Let  $S \subseteq \{1, 2, ..., N\} = [N]$ , and  $\vec{x} \in \mathbb{R}^N$ . Then, we let  $\vec{x}_S = \vec{x}$  with all entries not in S set to zero. That is,

$$(x_S)_j = \begin{cases} 0, & \text{if } j \neq S \\ x_j, & \text{if } j \in S \end{cases}$$

**Definition 4.** Similarly, if  $A \in \mathbb{R}^{m \times N}$ , then  $A_S :=$  the sub-matrix of A consisting of the columns of A indexed in S.

Example 2. 
$$\vec{x} = \begin{pmatrix} 100\\ 100\\ 100\\ 100 \end{pmatrix}$$
 and  $A$  is the 4 × 4 identity,  $A = I_{4\times4}$ , and  $S = \{2, 4\}$ , then  $\vec{x}_S = \begin{pmatrix} 0\\ 100\\ 0\\ 100 \end{pmatrix}$   
and  $A_S = \begin{pmatrix} 0 & 0\\ 1 & 0\\ 0 & 0\\ 0 & 1 \end{pmatrix}$ .

**Example 3.**  $\vec{x}_{supp(\vec{x})} = \vec{x} \quad \forall \vec{x} \in \mathbb{C}^N$ 

**Definition 5.** Let  $A \in \mathbb{C}^{m \times N}$ . Then  $k^{\text{th}}$  restricted isometry constant of A,  $\epsilon_k(A)$ , is

$$\epsilon_k(A) := \max_{S \subseteq [N], |S| \le k} \sigma_1 \left( A_S^* A_S - I_{|S| \times |S|} \right) = \max_{S \subseteq [N], |S| \le k} \max \left\{ \left| \sigma_1^2(A_S) - 1 \right|, \left| \sigma_{|S|}^2(A_S) - 1 \right| \right\}.$$

Note: If  $A \in \mathbb{R}^{m \times N}$  has the RIP of order k for  $\epsilon$ , then  $\epsilon_k(A) \leq \epsilon$ . Theorem 1 above is simply a means of bounding the restricted isometry constants of certain random matrices.

Any matrix with small RI constants will be useful for sparse approximation, as we will begin to see next.

**Lemma 1.** Let  $\vec{u}, \vec{v} \in \mathbb{C}^N$  be complex vectors s.t.  $\|\vec{u}\|_0 \leq k$ ,  $\|\vec{v}\|_0 \leq t$ . If  $supp(\vec{u}) \cap supp(\vec{v}) = \emptyset$ , then  $|\langle A\vec{u}, A\vec{v} \rangle| \leq \epsilon_{k+t}(A) \cdot \|\vec{u}\|_2 \cdot \|\vec{v}\|_2.$ 

*Proof:* Let  $S := supp(\vec{u}) \cup supp(\vec{v})$  so that |S| = t + k. Since  $\vec{u}$  and  $\vec{v}$  have disjoint supports, we have  $\langle \vec{u}, \vec{v} \rangle = 0$ . Thus,

$$\begin{split} | < A\vec{u}, A\vec{v} > | &= | < A_S \vec{u}_S, A_S \vec{v}_S > | \\ = | < A_S \vec{u}_S, A_S \vec{v}_S > - < \vec{u}_S, \vec{v}_S > | \\ = | < (A_S^* A_S - I) \vec{u}_S, \vec{v}_S > | \\ \le \| (A_S^* A_S - I) \| \cdot \| \vec{u}_S \|_2 \cdot \| \vec{v}_S \|_2 \quad \text{(by Hölder's Inequality)} \\ \le \epsilon_{k+t} (A) \cdot \| \vec{u}_S \|_2 \cdot \| \vec{v}_S \|_2. \end{split}$$

Note: If  $\vec{x} \in \mathbb{C}^N$ ,  $[N] = \bigcup_{j=0}^{\lfloor \frac{N}{k} \rfloor} S_j$  then  $\vec{x} = \sum_{j=0}^{\lfloor \frac{N}{k} \rfloor} \vec{x}_{S_j}$  where  $S_j \cap S_l = \emptyset$  if  $l \neq j$ , and  $|S_j| \leq k$  for all  $j = 0, 1, \ldots, \lfloor \frac{N}{k} \rfloor$ .

**Definition 6.** For any  $\vec{x} \in \mathbb{C}^N$ , we will let its coordinates be ordered by magnitude s.t.

$$|x_{i_1}| \ge |x_{i_2}| \ge \dots \ge |x_{i_N}|$$

and we will set  $S_0 = \{i_1, i_2, ..., i_k\}$ , and  $S_j = \{i_{1+jk}, i_{2+jk}, ..., i_{(j+1)k}\}$  for all  $j \in \lfloor \lfloor \frac{N}{k} \rfloor \rfloor$ . Lemma 2. If  $\vec{u} \in \mathbb{C}^N$  and  $\vec{V} \in \mathbb{C}^N$  satisfy

$$||u||_{\infty} := \max_{i \in [N]} |u_i| \le \min_{j \in [N]} |v_j|$$
(2)

then  $\|\vec{u}\|_2 \le \frac{\|\vec{v}\|_1}{\sqrt{N}}$ .

*Proof:* Homework 2.

**Theorem 2.** Suppose that  $\epsilon_{2k}(A) < \frac{1}{2}$  for some  $A \in \mathbb{C}^{m \times N}$ . Then

$$\|\vec{v}_S\|_1 \le \frac{\epsilon_{2k}(A)}{1 - 2\epsilon_{2k}(A)} \|\vec{v}_{\overline{S}}\|_1 \tag{3}$$

for all  $\vec{v} \in \ker(A)$  and  $\forall S \subseteq \{1, 2, 3, ..., N\}$  with  $|S| \leq k$ .

Note: (3) holding as above is commonly referred to as the "null space property" with constant  $\frac{\epsilon_{2k}(A)}{1-2\epsilon_{2k}(A)}$ . If  $A\vec{v} = 0$  then  $\vec{v}$  cannot be k-sparse.

*Proof:* Let  $S \subset [N]$  with  $|S| \leq k$ . If (3) holds for just  $S_0$  w.r.t. at given  $\vec{v}$  it is enough since

$$\|\vec{v}_S\|_1 \le \|\vec{v}_{S_0}\|_1 \le \frac{\epsilon_{2k}(A)}{1 - 2\epsilon_{2k}(A)} \|\vec{v}_{\overline{S_0}}\|_2 \le \frac{\epsilon_{2k}(A)}{1 - 2\epsilon_{2k}(A)} \|\vec{v}_{\overline{S}}\|_2 \tag{4}$$

Therefore, it suffices to establish (3) for  $S_0$  w.r.t. each  $\vec{v} \in \ker(A)$ 

By definition of the Restricted Isometry (RI) constant,  $\epsilon_{2k}(A) \geq \epsilon_k(A)$ . Thus

$$(1 - \epsilon_{2k}(A)) \|\vec{v}_{S_0}\|_2^2 \le (1 - \epsilon_k(A)) \|\vec{v}_{S_0}\|_2^2 \le \|A\vec{v}_{S_0}\|_2^2,$$
(5)

and so

$$\|v\vec{s}_0\|_2^2 \le \frac{1}{1 - \epsilon_{2k}(A)} \|A\vec{v}_{S_0}\|_2^2 = \frac{1}{1 - \epsilon_{2k}(A)} \langle A\vec{v}_{S_0}, A\vec{v}_{S_0} \rangle.$$

Since  $\vec{v} \in \ker(A)$ , we have that  $A\vec{v}_{S_0} = -\sum_{j=1}^{\lfloor \frac{N}{k} \rfloor} Av_{S_j}$ . Thus, Lemma 1 implies that

$$\|\vec{v}_{S_0}\|_2^2 \le \frac{1}{1 - \epsilon_{2k}(A)} \sum_{j \ge 1} \langle A\vec{v}_{S_0}, -A\vec{v}_{S_j} \rangle \le \frac{\epsilon_{2k}(A)}{1 - \epsilon_{2k}(A)} \sum_{j \ge 1} \|\vec{v}_{S_0}\|_2 \|\vec{v}_{S_j}\|_2$$

and so we get that

$$\|\vec{v}_{S_0}\| \le \frac{\epsilon_{2k}(A)}{1 - \epsilon_{2k}(A)} \sum_{j=1}^{\lfloor \frac{N}{k} \rfloor} \|\vec{v}_{S_j}\|_2.$$
(6)

Now, Lemma 2 tells us that  $\|\vec{v}_{S_j}\|_2 \leq \frac{\|\vec{v}_{S_{j-1}}\|_1}{\sqrt{k}}, \ \forall j \geq 1$ . Thus,

$$\|\vec{v}_{S_0}\|_2 \le \frac{1}{\sqrt{k}} \cdot \frac{\epsilon_{2k}(A)}{1 - \epsilon_{2k}(A)} \cdot \|\vec{v}\|_1.$$

Now to finish, we note that Hölder's Inequality implies that  $\|\vec{v}_{S_0}\|_1 \leq \|\vec{v}_{S_0}\|_2 \cdot \sqrt{k}$ . Therefore, a rearrangement of the last inequality will give us the desired result.

## 1.1 Homework Problems:- due March 21<sup>st</sup> (Tue.)

Homework 1: Prove theorem 1. Homework 2: Prove lemma 2.

## References

[1] Simon Foucart and Holger Rauhut. A Mathematical Introduction to Compressive Sensing. Birkhauser Basel, 2013.