NAME:

## Math 481

## Quiz 24 Solution

**1.** Define the distance between vertices in a connected graph: dist(v, w) means ...?

Solution: For vertices v, w in a graph G, the distance dist(v, w) is the minimum number of edges in a path from v to w inside G, i.e. the length of a shortest path  $vPw \subset G$ .

*Note:* We could also say "minumum number of edges in a *walk*", since the shortest walk is always a path (no repeat vertices).

**2.** TRUE OR FALSE: For any distinct vertices x, y, z in a connected graph,

$$\operatorname{dist}(x, z) = \operatorname{dist}(x, y) + \operatorname{dist}(y, z).$$

Solution: False. A counterexample is the complete graph  $K_3$  with vertices  $V = \{x, y, z\}$ , having dist(x, y) = dist(x, z) = dist(y, z) = 1, but  $1 \neq 1 + 1$ .

**3.** TRUE OR FALSE: For any distinct vertices x, y, z in a connected graph,

$$\operatorname{dist}(x, z) \le \operatorname{dist}(x, y) + \operatorname{dist}(y, z).$$

Solution: True. Informally, a shortest path from x to y joined to a shortest path from y to z is a walk from x to z with dist(x, y) + dist(y, z) edges, and the shortest xz-walk (or path), having dist(x, z) edges, cannot be longer than this.

Formal proof: Let xPy and yQz be shortest paths, by definition having length dist(x, y) and dist(y, z). Then their concatenation xPyQz is a walk from x to z of length dist(x, y) + dist(y, z). This is not necessarily a path, since Q might cross or backtrack P; however, we know that any walk can be cut down to a path with the same endpoints by skipping any loops. Thus there exists a path from x to z of length at most dist(x, y) + dist(y, z), and a shortest path of length dist(x, z) cannot be longer than this, which is the conclusion to be proved.

*Note:* This remains true for a disconnected graph, in which some distances can be infinite: if  $dist(x, z) = \infty$ , then dist(x, y) or dist(y, z) must be  $\infty$ .