

**Basic number sets**

Natural numbers  $\mathbb{N} = \{0, 1, 2, 3, 4, 5, \dots\}$ , infinite set, ordering  $a < b$ , Well-Ordering Axiom

Integers  $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$  with operations  $+, -, \times$

Rational numbers  $\mathbb{Q} = \{\frac{a}{b} \mid a \in \mathbb{Z}, b \in \mathbb{N}\}$  with operations  $+, -, \times, \div$

**Division with remainder**

**THEOREM 1.1** For  $a, b \in \mathbb{Z}$ ,  $b > 0$ , there exist unique  $q, r \in \mathbb{Z}$  such that  $a = qb + r$  and  $0 \leq r < b$ .

*Proof of Existence.*

$$\begin{aligned} \mathbb{Z} = & \dots \cup \{-b, -b+1, \dots, -1\} \cup \{0, 1, \dots, b-1\} \\ & \cup \{b, b+1, \dots, 2b-1\} \cup \{2b, 2b+1, \dots, 3b-1\} \cup \dots \end{aligned}$$

Then  $a$  lies in one of the above intervals,  $a \in \{nb, nb+1, \dots, nb+i, \dots, nb+b-1\}$  for some  $n \in \mathbb{Z}$ , which means  $a = nb + i$  with  $0 \leq i < b$ . Thus  $q = n$ ,  $r = i$  satisfies the required properties.

*Proof of Uniqueness.* Suppose  $a = qb + r = q'b + r'$  with  $0 \leq r, r' < b$ . Then:

$$r - r' = (a - qb) - (a - q'b) = (q - q')b,$$

which is a multiple of  $b > 0$ . But:

$$-b < -r' \leq r - r' \leq r < b,$$

and the only multiple of  $b$  in this interval is zero, so  $r - r' = (q - q')b = 0$ . Thus  $r - r' = 0$  and  $q - q' = 0$ , so  $r = r'$  and  $q = q'$ . ■

Long division algorithm  $b \overline{)a}$  gives method for computing  $q, r$ .

**Euclidean Algorithm.**

For  $a, b \in \mathbb{N}$ , their *greatest common divisor*  $d = \gcd(a, b) = (a, b)$  is the largest integer dividing both  $a$  and  $b$  (with no remainder).

**EXAMPLE:** The divisors of 12 are 1, 2, 3, 4, 6, 12; divisors of 8 are 1, 2, 4, 8. Thus  $\gcd(12, 8) = 4$ .

The Euclidean Algorithm is an efficient method to find  $\gcd(a, b)$  for  $a \geq b > 0$  by repeated division with remainder:

$$\begin{aligned} a &= q_1 b + r_1 \\ b &= q_2 r_1 + r_2 \\ r_1 &= q_3 r_2 + r_3 \\ &\vdots \\ r_{n-2} &= q_n r_{n-1} + r_n \\ r_{n-1} &= q_{n+1} r_n + 0. \end{aligned}$$

This process must terminate since  $a \geq b > r_1 > r_2 > \dots \geq 0$ , and we cannot keep decreasing infinitely by Well-Ordering.

Formally, if we start with  $r_{-1} = a$ ,  $r_0 = b$ , and we have already computed  $r_1, \dots, r_{i-1}$  for some  $i \geq 1$ , then  $q_i, r_i$  are computed recursively by the formula:

$$r_{i-2} = q_i r_{i-1} + r_i,$$

stopping when  $r_i = 0$ . Next time, we will prove that the gcd is the last non-zero remainder:

$$r_n \stackrel{!!}{=} \gcd(a, b).$$

*Example.* Find  $\gcd(57, 21)$ .

$$\begin{aligned} 57 &= 2 \cdot 21 + 15 \\ 21 &= 1 \cdot 15 + 6 \\ 15 &= 2 \cdot 6 + 3 \\ 6 &= 2 \cdot 3 + 0. \end{aligned}$$

Here  $a = 57 > b = 21 > r_1 = 15 > r_2 = 6 > r_3 = 3 > r_4 = 0$ .

The gcd is the last non-zero remainder:  $\gcd(57, 21) = r_3 = 3$ .

### Extended Euclidean Algorithm

For  $\gcd(a, b) = d$ , we can find integers  $s, t \in \mathbb{Z}$  such that

$$d = sa + tb,$$

by starting with  $d = r_n = r_{n-2} - q_n r_{n-1}$  and successively substituting  $r_i = r_{i-2} - q_i r_{i-1}$  for  $i = n-1, n-2, \dots, 1$ .

In our example:

$$\begin{aligned} d = 3 &= 15 - 2 \cdot 6 \\ 6 &= 21 - 1 \cdot 15 \\ 15 &= 57 - 2 \cdot 21. \end{aligned}$$

Successively substituting:

$$\begin{aligned} d = 3 &= 15 - 2 \cdot 6 \\ &= 15 - 2(21 - 1 \cdot 15) &= -2 \cdot 21 + 3 \cdot 15 \\ &= -2 \cdot 21 + 3 \cdot (57 - 2 \cdot 21) &= 3 \cdot 57 - 8 \cdot 21, \end{aligned}$$

so we can take  $s = 3$ ,  $t = -8$  to satisfy  $3 = s \cdot 57 + t \cdot 21$ .