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Proposition: Let $G$ be a connected graph containing a cycle $C$ with edge $e=x y$, so $e \subset C \subset G$. Then the graph $G-e$, removing $e$ but keeping vertices $x, y$, is connected.

Sketch a proof, first writing the hypothesis at the top and the conclusion at the bottom, with their definitions. Notation: let $C=x e y Q x$, from $x$ along $e$ to $y$, then along path $Q$ back to $x$.

Informal argument: We need to show any two vertices $v, w$ are connected in $G-e$. There is a path $v P w$ in $G$, which might contain the edge $e=x y$ in $C$. In that case, avoid $e$ by going the other way around $C$, getting a new path which connects $v$ to $w$ in $G-e$.

Formal proof: By hypothesis, assume that $G$ is a connected graph, meaning any vertices $v, w$ are connected by a path $v P w \subset G$. Also, $G$ contains a cycle of the form $C=x e y Q x$, where $e=x y$ is an edge and $y Q x$ is a path not containing $e$. Let $G-e$ be the graph with the same vertices as $G$ but missing edge $e$.

Since $G$ is connected, any two vertices $v, w$ are connected by some path $v P w$ in $G$. If $P$ does not contain $e$, then $P^{\prime}=P$ is a path $v P^{\prime} w \subset G-e$.

If $P$ does contain $e=x y$, we may write it as $v P_{1} x e y P_{2} w$, and construct

$$
R=v P_{1} x Q y P_{2} w
$$

the walk from $v$ along $P$ to $x$, then around $Q$ to $y$, continuing along $P$ to $w$. By definition, the path $P$ and the cycle $C$ can use $e$ only once, so the path segments $v P x, x Q y, y P w$ cannot contain $e$. Thus $R$ is a walk in $G-e$ connecting $v, w$, and we can cut it down to a path $v P^{\prime} w \subset G-e$.

Either way, arbitrary vertices $v, w$ are connected by $v P^{\prime} w$ in $G-e$, so it is a connected graph.

