

2e. Mycielski gave an iterative construction of a triangle-free graph of arbitrarily high chromatic number. Given a triangle-free G with n vertices and $\chi(G) = k$, we define a new triangle-free graph M with $2n+1$ vertices and $\chi(M) = k+1$.

Let M contain a copy of G with vertices $\{v_1, \dots, v_n\}$, as well as a star graph $K_{1,n}$ with vertices $\{u_0, u_1, \dots, u_n\}$ and edges u_0u_i . Keep each edge v_iv_j in G , and add two edges v_iv_j, v_ju_i in M .

The only possible triangle in M is of the form $v_iv_ju_k$, but this would occur only if $v_iv_jv_k$ were a triangle in G .

To see that the construction increases the chromatic number $\chi(G) = k$, consider a proper k -coloring of $M - \{u_0\}$; that is, a mapping $c : \{v_1, \dots, v_n, u_1, \dots, u_n\} \rightarrow \{1, 2, \dots, k\}$ with $c(x) \neq c(y)$ for adjacent vertices x, y . If we had $c(u_i) \in \{1, 2, \dots, k-1\}$ for all i , then we could define a proper $(k-1)$ -coloring of G by $c'(v_i) = c(u_i)$ when $c(v_i) = k$, and $c'(v_i) = c(v_i)$ otherwise. But this is impossible for $\chi(G) = k$, so c must use all k colors for $\{u_1, \dots, u_n\}$, and any proper coloring of the last vertex u_0 must use an extra color. That is, $\chi(M) = k+1$.

3. PROPOSITION: For fixed k and $p(n) = \frac{\log(n)}{\sqrt{n}}$, a graph $G \in \mathcal{G}(n, p(n))$ is almost certainly k -connected (in fact, connected by 2-edge paths).

Proof: To show $\mathbb{P}(G \in \mathcal{G}(n, p(n)) \text{ is } k\text{-connected}) \rightarrow 1$ as $n \rightarrow \infty$, we will bound the probability that G fails to be k -connected. Writing $G = (V, E)$ with $|V| = n$ vertices, we ask if there is a k -element vertex set $S \subset V$ for which $G - S$ is not connected by 2-edge paths, so that there are $a, b \in V - S$ such that no path acb lies in G , for $c \in V - S - \{a, b\}$:

$$\begin{aligned} \mathbb{P}(\text{FAIL}) &\leq \mathbb{P}(\exists S \exists a, b \forall c : ac \notin E \text{ or } bc \notin E) \\ &\leq \sum_S \sum_{a, b} \mathbb{P}(\forall c : \text{NOT}(ac, bc \in E)). \end{aligned}$$

Now, for $c \neq c'$, the edges ac, bc, ac', bc' are distinct, so the events $ac, bc \in G$ and $ac', bc' \in G$ are independent. Hence:

$$\begin{aligned} \mathbb{P}(\text{FAIL}) &\leq \sum_S \sum_{a, b} \prod_c (1 - \mathbb{P}(ac, bc \in G)) \\ &= \binom{n}{k} \binom{n-k}{2} (1-p^2)^{n-k-2} \\ &\leq An^{k+2}(1-p^2)^n = An^{k+2} \left(1 - \frac{\log^2(n)}{n}\right)^n, \end{aligned}$$

where $A = A(k)$ is a constant. We must show this upper bound tends to zero as $n \rightarrow \infty$.

Taking the logarithm and using $\log(1-x) \leq -x$ gives:

$$\log(A) + (k+2) \log(n) + n \log\left(1 - \frac{\log^2(n)}{n}\right) \leq \log(A) + (k+2) \log(n) - n \left(\frac{\log^2(n)}{n}\right).$$

Since $\log^2(n) \gg (k+2) \log(n)$, it is clear that $\lim_{n \rightarrow \infty} \log(An^{k+2}(1-p^2)^n) = -\infty$, and:

$$\mathbb{P}(\text{FAIL}) \leq An^{k+2}(1-p^2)^n \rightarrow 0 \text{ as } n \rightarrow \infty.$$

Hence, the complementary probability that G is k -connected tends to 1.

Note that this argument works for any $p(n) = \frac{\sqrt{\log(n)}}{\sqrt{n}} f(n)$ with $f(n) \rightarrow \infty$.