

Fast Robust Phase Retrieval from Local Correlation Measurements

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The Phase Retrieval Problem

$$\text{find}^1 \quad \mathbf{x} \in \mathbb{C}^d \quad \text{given} \quad |M\mathbf{x}|^2 + \mathbf{n} = \mathbf{b} \in \mathbb{R}^D,$$

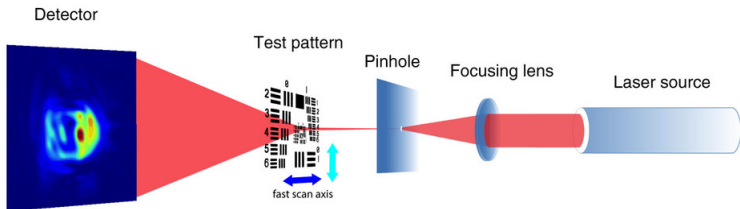
where

- $\mathbf{b} \in \mathbb{R}^D$ denotes the phaseless (or magnitude-only) measurements,
- $M \in \mathbb{C}^{D \times d}$ is a measurement matrix associated with these measurements, and
- $\mathbf{n} \in \mathbb{R}^D$ is measurement noise.

Let $\mathcal{A} : \mathbb{R}^D \rightarrow \mathbb{C}^d$ denote the recovery method. The phase retrieval problem involves designing measurement matrix and recovery method pairs, (M, \mathcal{A}) .

¹(upto a global phase offset)

Motivating Applications



From Huang, Xiaojing, et al. "Fly-scan ptychography." *Scientific Reports* 5 (2015).

The Phase Retrieval problem arises in many molecular imaging modalities, including

- X-ray crystallography
- Ptychography

Other applications can be found in optics, astronomy and speech processing.

Existing Computational Approaches

- Alternating projection methods
[Fienup, 1978], [Fannjiang and Liao, 2012], ...
- Methods based on semidefinite programming
PhaseLift [Candes et al., 2012], PhaseCut [Waldspurger et al., 2012], ...
- Others
 - Frame-theoretic, graph based algorithms [Alexeev et al., 2014]
 - (Stochastic) gradient descent [Candes et al., 2014]

... and variants for sparse and/or structured signal models.

Existing Computational Approaches

- Alternating projection methods – **No recovery guarantees** [Fienup, 1978], [Fannjiang and Liao, 2012], ...
- Methods based on semidefinite programming – **Expensive** PhaseLift [Candes et al., 2012], PhaseCut [Waldspurger et al., 2012], ...
- Others
 - Frame-theoretic, graph based algorithms [Alexeev et al., 2014]
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Most methods require **random** measurement constructions.

Today...

- We discuss a recently introduced **essentially linear-time** phase retrieval algorithm based on (**deterministic²**) **local correlation measurement constructions**.
- We provide rigorous theoretical recovery guarantees and present numerical results showing the accuracy, efficiency and robustness of the method.

²for a large class of real-world signals

Outline

- 1 The Phase Retrieval Problem
- 2 BlockPR: Fast Phase Retrieval from Local Correlation Measurements
 - Measurement Constructions
 - Solving for Phase Differences
 - Angular Synchronization
- 3 Theoretical Guarantees
- 4 Numerical Simulations

Key Components

- 1 **Local Measurements:** Each measurement provides information about some *local* region of \mathbf{x} .
- 2 **Local Lifting:** Use compactly supported masks and correlation measurements to obtain phase difference estimates.

$$|\text{Corr}(\mathbf{m}_i, \mathbf{x})|^2 \xrightarrow[\text{linear system}]{\text{solve}} \{x_j x_k^*\}_{|j-k| \bmod d < \delta}$$

- \mathbf{m}_i is a mask or window function with δ non-zero entries.
- $x_j x_k^*$ provides (scaled) phase difference between x_j and x_k .

- 3 **Angular Synchronization:** Use the phase differences to obtain the phases of the unknown signal.

$$\{x_j x_k^*\}_{|j-k| \bmod d < \delta} \xrightarrow[\text{synchronization}]{\text{angular}} \{x_j\}_{j=1}^d$$

Constraint on \mathbf{x} : We require \mathbf{x} to be “flat”.

(At most δ consecutive entries in \mathbf{x} with magnitude $< \frac{\|\mathbf{x}\|_2}{2\sqrt{d}}$)

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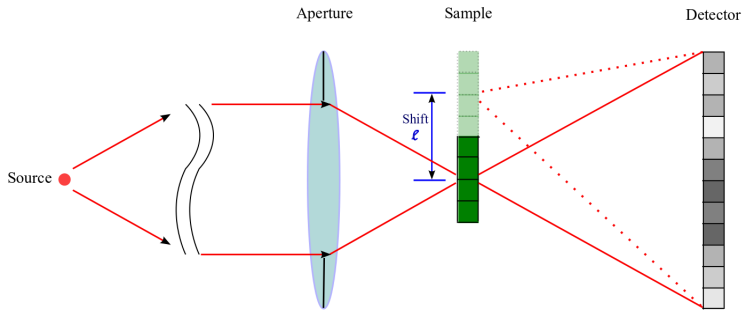
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Measurement Constructions



Adapted from Huang et al. "Fly-scan ptychography." *Scientific Reports* 5 (2015).

- We consider measurements motivated by **Ptychographic** molecular imaging.
- Measurements are **local**; the full reconstruction is obtained by imaging *shifts* of the specimen.

Model Problem

Recover an unknown vector $\mathbf{x} \in \mathbb{C}^4$ from noiseless measurements

$$\mathbf{y} = |M\mathbf{x}|^2,$$

where $\mathbf{y} \in \mathbb{R}^{12}$ and $M \in \mathbb{C}^{12 \times 4}$ has the following structure:

$$M = \begin{bmatrix} M_1 \\ M_2 \\ M_3 \end{bmatrix}, \quad M_i = \begin{bmatrix} (\mathbf{m}_i)_1 & (\mathbf{m}_i)_2 & 0 & 0 \\ 0 & (\mathbf{m}_i)_1 & (\mathbf{m}_i)_2 & 0 \\ 0 & 0 & (\mathbf{m}_i)_1 & (\mathbf{m}_i)_2 \\ (\mathbf{m}_i)_2 & 0 & 0 & (\mathbf{m}_i)_1 \end{bmatrix}.$$

Here, $\mathbf{m}_{\{1,2,3\}} \in \mathbb{C}^4$ are masks with *local support* (with $\delta = 2$ non-zero entries).

Local Correlation Measurements

These correspond to *local correlation measurements*

$$\begin{aligned} (\mathbf{y}_\ell)_i &= \left| \sum_{k=1}^{\delta=2} (\mathbf{m}_\ell)_k \cdot x_{i+k-1} \right|^2, & (\ell, i) &\in \{1, 2, 3\} \times \{1, 2, 3, 4\} \\ &= \sum_{j,k=1}^{\delta} (\mathbf{m}_\ell)_j (\mathbf{m}_\ell)_k^* x_{i+j-1} x_{i+k-1}^* := \sum_{j,k=1}^{\delta} (\mathbf{m}_\ell)_{j,k} x_{i+j-1} x_{i+k-1}^*. \end{aligned}$$

This is a **linear** system for the phase differences $\{x_j x_k^*\}$!

Note: the masks $\mathbf{m}_{\{1,2,3\}}$ (which are related to the aperture transmission function of the imaging system) are known - either by design or through calibration.

Solving for Phase Differences

Writing out the correlation sum, we obtain the linear system

$$M' \mathbf{z} = \tilde{\mathbf{b}},$$

where

$$\mathbf{z} = [|x_1|^2 \quad x_1 x_2^* \quad x_2 x_1^* \quad |x_2|^2 \quad x_2 x_3^* \quad x_3 x_2^* \quad |x_3|^2 \quad x_3 x_4^* \quad x_4 x_3^* \quad |x_4|^2 \quad x_4 x_1^* \quad x_1 x_4^*]^T,$$

$$\tilde{\mathbf{b}} = [(y_1)_1 \quad (y_2)_1 \quad (y_3)_1 \quad (y_1)_2 \quad (y_2)_2 \quad (y_3)_2 \quad (y_1)_3 \quad (y_2)_3 \quad (y_3)_3 \quad (y_1)_4 \quad (y_2)_4 \quad (y_3)_4]^T,$$

$$M' = \begin{bmatrix} (\mathbf{m}_1)_{1,1} & (\mathbf{m}_1)_{1,2} & (\mathbf{m}_1)_{2,1} & (\mathbf{m}_1)_{2,2} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ (\mathbf{m}_2)_{1,1} & (\mathbf{m}_2)_{1,2} & (\mathbf{m}_2)_{2,1} & (\mathbf{m}_2)_{2,2} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ (\mathbf{m}_3)_{1,1} & (\mathbf{m}_3)_{1,2} & (\mathbf{m}_3)_{2,1} & (\mathbf{m}_3)_{2,2} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & (\mathbf{m}_1)_{1,1} & (\mathbf{m}_1)_{1,2} & (\mathbf{m}_1)_{2,1} & (\mathbf{m}_1)_{2,2} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & (\mathbf{m}_2)_{1,1} & (\mathbf{m}_2)_{1,2} & (\mathbf{m}_2)_{2,1} & (\mathbf{m}_2)_{2,2} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & (\mathbf{m}_3)_{1,1} & (\mathbf{m}_3)_{1,2} & (\mathbf{m}_3)_{2,1} & (\mathbf{m}_3)_{2,2} & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & (\mathbf{m}_1)_{1,1} & (\mathbf{m}_1)_{1,2} & (\mathbf{m}_1)_{2,1} & (\mathbf{m}_1)_{2,2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & (\mathbf{m}_2)_{1,1} & (\mathbf{m}_2)_{1,2} & (\mathbf{m}_2)_{2,1} & (\mathbf{m}_2)_{2,2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & (\mathbf{m}_3)_{1,1} & (\mathbf{m}_3)_{1,2} & (\mathbf{m}_3)_{2,1} & (\mathbf{m}_3)_{2,2} & 0 & 0 \\ (\mathbf{m}_1)_{2,2} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & (\mathbf{m}_1)_{1,1} & (\mathbf{m}_1)_{1,2} & (\mathbf{m}_1)_{2,1} \\ (\mathbf{m}_2)_{2,2} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & (\mathbf{m}_2)_{1,1} & (\mathbf{m}_2)_{1,2} & (\mathbf{m}_2)_{2,1} \\ (\mathbf{m}_3)_{2,2} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & (\mathbf{m}_3)_{1,1} & (\mathbf{m}_3)_{1,2} & (\mathbf{m}_3)_{2,1} \end{bmatrix}$$

Back to our Example ...

$$[|x_1|^2 \quad x_1x_2^* \quad x_2x_1^* \quad |x_2|^2 \quad x_2x_3^* \quad x_3x_2^* \quad |x_3|^2 \quad x_3x_4^* \quad x_4x_3^* \quad |x_4|^2 \quad x_4x_1^* \quad x_1x_4^*]^T$$

↓ (re-arrange)

$$\begin{bmatrix} |x_1|^2 & x_1x_2^* & 0 & x_1x_4^* \\ x_2x_1^* & |x_2|^2 & x_2x_3^* & 0 \\ 0 & x_3x_2^* & |x_3|^2 & x_3x_4^* \\ x_4x_1^* & 0 & x_4x_3^* & |x_4|^2 \end{bmatrix} \quad (2\delta - 1 \text{ entries in band})$$

↓ (normalize)

$$\begin{bmatrix} 1 & e^{i(\phi_1 - \phi_2)} & 0 & e^{i(\phi_1 - \phi_4)} \\ e^{i(\phi_2 - \phi_1)} & 1 & e^{i(\phi_2 - \phi_3)} & 0 \\ 0 & e^{i(\phi_3 - \phi_2)} & 1 & e^{i(\phi_3 - \phi_4)} \\ e^{i(\phi_4 - \phi_1)} & 0 & e^{i(\phi_4 - \phi_3)} & 1 \end{bmatrix}$$

↓ (angular synchronization)

$$\phi_1, \phi_2, \phi_3, \phi_4$$

(Signal Reconstruction) $[|x_1|e^{i\phi_1} \quad |x_2|e^{i\phi_2} \quad |x_3|e^{i\phi_3} \quad |x_4|e^{i\phi_4}]^T$

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The Angular Synchronization Problem

The Angular Synchronization Problem

Estimate d unknown angles $\phi_1, \phi_2, \dots, \phi_d \in [0, 2\pi)$ from noisy and possibly incomplete measurements of their differences,

$$\phi_{i,j} := \phi_i - \phi_j \pmod{2\pi}.$$

- Several possible approaches: eigenvector methods, semidefinite programming . . .
- Today: Greedy angular synchronization

Greedy Angular Synchronization

- 1 Set the largest magnitude component to have zero phase angle; i.e.,

$$\arg(x_j) = 0, \quad j = \operatorname{argmax}_i |x_i|^2.$$

- 2 Use this entry to set the phase angles of its δ neighboring entries; i.e.,

$$\arg(x_k) = \arg(x_j) - \phi_{j,k}, \quad |j - k \bmod d| < \delta.$$

- 3 Use the next largest magnitude component from these δ entries and repeat the process.

Greedy Angular Synchronization

$$[|x_1|^2 \quad x_1 x_2^* \quad x_2 x_1^* \quad |x_2|^2 \quad x_2 x_3^* \quad x_3 x_2^* \quad |x_3|^2 \quad x_3 x_4^* \quad x_4 x_3^* \quad |x_4|^2 \quad x_4 x_1^* \quad x_1 x_4^*]^T$$

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Greedy Angular Synchronization

Applying this to our example problem. . .

- Assume, without loss of generality, that $|x_1| \geq |x_i|$, $i \in \{2, 3, 4\}$.
- ① We start by setting³ $\arg(x_1) = 0$.
- ② We may now set the phase of x_2 and x_4 using the estimated phase differences $\phi_{1,2}$ and $\phi_{1,4}$ respectively; i.e.,

$$\arg(x_2) = \arg(x_1) - \phi_{1,2}, \quad \arg(x_4) = \arg(x_1) - \phi_{1,4}.$$

- ③ Similarly, we next set $\arg(x_3) = \arg(x_2) - \phi_{2,3}$, thereby recovering all of the entries' unknown phases.

³Recall that we can only recover \mathbf{x} up to an unknown global phase factor which, in this case, will be the true phase of x_1 .

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Block-Circulant Matrix: Condition Number Bounds

Theorem (Iwen, V., Wang 2015)

Choose entries of the measurement mask (\mathbf{m}_i) as follows:

$$(\mathbf{m}_i)_\ell = \begin{cases} \frac{e^{-\ell/a}}{\sqrt[4]{2\delta-1}} \cdot e^{\frac{2\pi i \cdot i \cdot \ell}{2\delta-1}}, & \ell \leq \delta \\ 0, & \ell > \delta \end{cases}, \quad \begin{matrix} a := \max \left\{ 4, \frac{\delta-1}{2} \right\}, \\ i = 1, 2, \dots, N. \end{matrix}$$

Then, the resulting system matrix for the phase differences, M' , has condition number

$$\kappa(M') < \max \left\{ 144e^2, \frac{9e^2}{4} \cdot (\delta - 1)^2 \right\}.$$

- **Deterministic** (windowed DFT-type) measurement masks!
- δ is typically chosen to be $c \log_2 d$ with c small (2–3).
- Extensions: oversampling, random masks

Recovery Guarantee – Non-Sparse (“Flat”) Signals

Theorem (Iwen, V., Wang 2015)

There exist fixed universal constants $C, C' \in \mathbb{R}^+$ such that following holds: Let $M \in \mathbb{C}^{D \times d}$ be defined as in the previous slide, and suppose that $\mathbf{x} \in \mathbb{C}^d$ is non-sparse^a with $d > 2$ and $\|\mathbf{x}\|_2^2 \geq C(\delta - 1)d^2 \|\mathbf{n}\|_2$. Then, the proposed algorithm is guaranteed to recover an $\tilde{\mathbf{x}} \in \mathbb{C}^d$ with

$$\min_{\theta \in [0, 2\pi)} \left\| \mathbf{x} - e^{i\theta} \tilde{\mathbf{x}} \right\|_2^2 \leq C' d^2 (\delta - 1) \|\mathbf{n}\|_2$$

when given arbitrarily noisy input measurements

$\mathbf{b} = |M\mathbf{x}|^2 + \mathbf{n} \in \mathbb{R}^D$. Furthermore, the algorithm requires just $\mathcal{O}(\delta \cdot d \log d)$ operations for this choice of $M \in \mathbb{C}^{D \times d}$.

^adoes not have more than $\lfloor (\delta - 3)/2 \rfloor$ consecutive zeros or small entries; see preprint for details.

Recovery Guarantee - Arbitrary Signals

Theorem (Iwen, V., Wang 2015)

Let $\mathbf{x} \in \mathbb{C}^d$ with d sufficiently large have $\|\mathbf{x}\|_2^2 \geq C (d \ln d)^2 \ln^3(\ln d) \|\mathbf{n}\|_2$.^a Then, one can select a random measurement matrix $\tilde{M} \in \mathbb{C}^{D \times d}$ such that the following holds with probability at least $1 - \frac{1}{C' \cdot \ln^2(d) \cdot \ln^3(\ln d)}$: the proposed algorithm will recover an $\tilde{\mathbf{x}} \in \mathbb{C}^d$ with

$$\min_{\theta \in [0, 2\pi)} \left\| \mathbf{x} - e^{i\theta} \tilde{\mathbf{x}} \right\|_2^2 \leq C'' (d \ln d)^2 \ln^3(\ln d) \|\mathbf{n}\|_2$$

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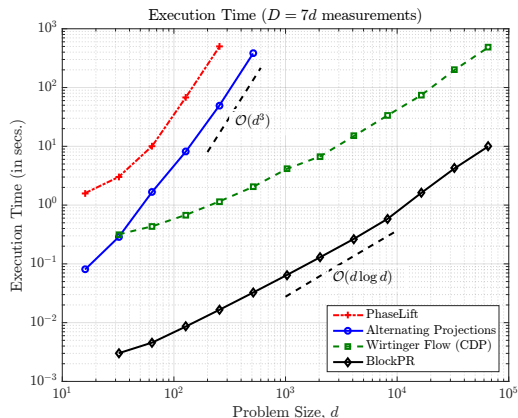
$\mathbf{b} = |\tilde{M}\mathbf{x}|^2 + \mathbf{n} \in \mathbb{R}^D$. Here D can be chosen to be $\mathcal{O}(d \cdot \ln^2(d) \cdot \ln^3(\ln d))$. Furthermore, the algorithm will run in $\mathcal{O}(d \cdot \ln^3(d) \cdot \ln^3(\ln d))$ -time.

^aHerein $C, C', C'' \in \mathbb{R}^+$ are all fixed and absolute constants.

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- 3 Theoretical Guarantees
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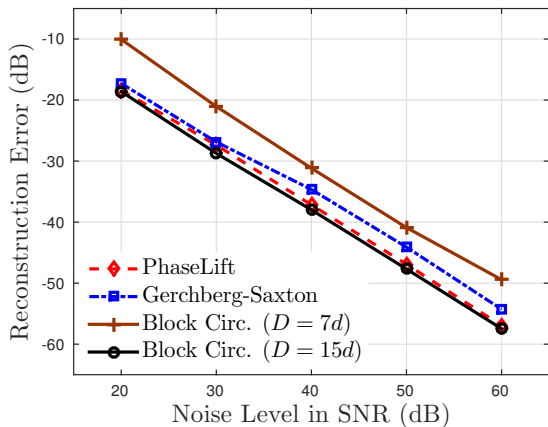
Efficiency



- iid Complex Gaussian test signal
- Averaged over 100 trials
- Simulations performed in Matlab on a laptop computer with 4GB RAM

Robustness

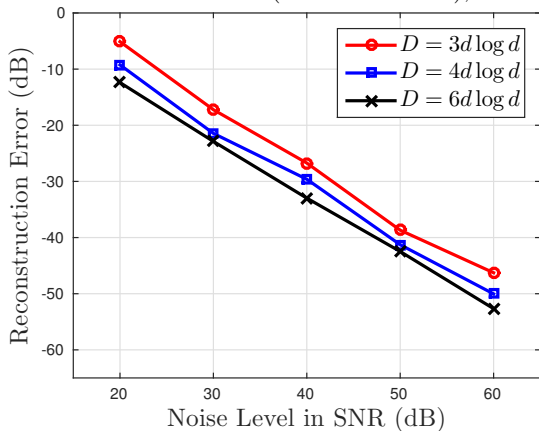
Robustness to Additive Noise, $d = 64, D = 7d$



- iid complex Gaussian signal
- $d = 64$
- $7d$ measurements
- Deterministic (windowed Fourier-like) measurements

Robustness

Robustness to Noise (Random Masks), $d = 2048$



- iid complex Gaussian signal
- $d = 2048$
- Not feasible with SDP-based methods such as PhaseLift on a laptop in Matlab
- Random measurements

In Summary. . .

- BlockPR allows for **essentially linear-time** robust phase retrieval from **local correlation measurement constructions**.
- **Deterministic** measurements for flat vectors.
- **First** known global *robust* recovery guarantee for phase retrieval from local correlation (ptychographic) measurements.

Current and Future Directions

- (Sublinear-time) compressive phase retrieval
- Improved angular synchronization frameworks
- Extensions to 2D and Ptychography

Publications/Preprints/Code

This Talk

Mark Iwen, A. Viswanathan and Yang Wang. “Fast Phase Retrieval from Local Correlation Measurements.” arXiv:1501.02377, 2015.

Code: <https://bitbucket.org/charms/blockpr>

Related Work

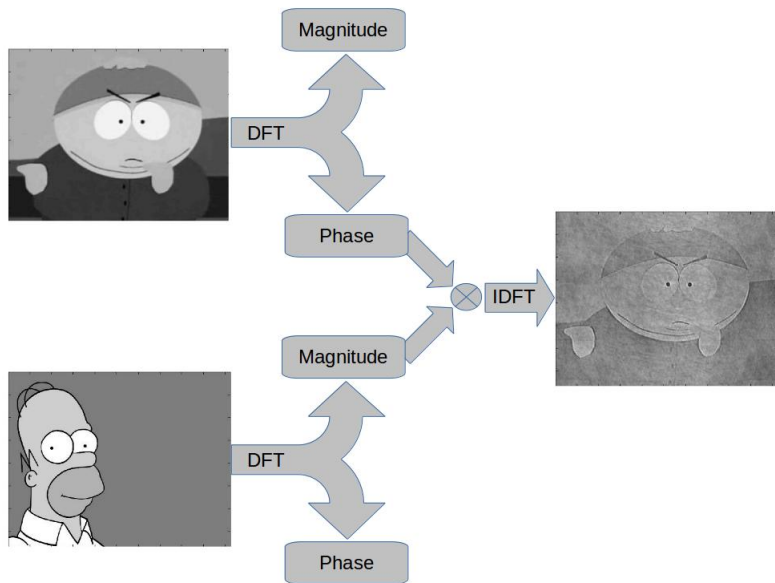
M. Iwen, A. Viswanathan, and Y. Wang. “Robust Sparse Phase Retrieval Made Easy.” (in press) ACHA, 2015. arXiv:1410.5295

A. Viswanathan and Mark Iwen. “Fast Angular Synchronization for Phase Retrieval via Incomplete Information.” Proc. SPIE 9597, Wavelets+Sparsity XVI, 2015.

A. Viswanathan and Mark Iwen. “Fast Compressive Phase Retrieval.” Asilomar Conf. Signals, Systems Computers, 2015.

Code: <https://bitbucket.org/charms/sparsepr>

Questions?



Appendix: Condition Number Proof Sketch

(Step 1) M' is block-circulant and therefore admits a unitary decomposition

$$U_{2\delta-1}^* M' U_{2\delta-1} = J = \text{blockdiag} (J_1, J_2, \dots, J_d),$$

where $J_1, \dots, J_d \in \mathbb{C}^{(2\delta-1) \times (2\delta-1)}$ are defined as

$$J_k := \sum_{l=1}^{\delta} M'_l \cdot e^{\frac{2\pi i \cdot (k-1) \cdot (l-1)}{d}}$$

and $U_\alpha \in \mathbb{C}^{\alpha d \times \alpha d}$ are unitary block Fourier matrices defined by

$$U_\alpha := \frac{1}{\sqrt{d}} \begin{pmatrix} I_\alpha & I_\alpha & \dots & I_\alpha \\ I_\alpha & I_\alpha e^{\frac{2\pi i}{d}} & \dots & I_\alpha e^{\frac{2\pi i \cdot (d-1)}{d}} \\ & & \ddots & \\ I_\alpha & I_\alpha e^{\frac{2\pi i \cdot (d-2)}{d}} & \dots & I_\alpha e^{\frac{2\pi i \cdot (d-2) \cdot (d-1)}{d}} \\ I_\alpha & I_\alpha e^{\frac{2\pi i \cdot (d-1)}{d}} & \dots & I_\alpha e^{\frac{2\pi i \cdot (d-1) \cdot (d-1)}{d}} \end{pmatrix}.$$

Appendix: Condition Number Proof Sketch

(Step 2) For the prescribed structured measurements, evaluating J_k yields

$$J_k = F_{2\delta-1} \begin{pmatrix} s_{k,1} & 0 & \dots & 0 \\ 0 & s_{k,2} & 0 & \dots \\ 0 & 0 & \ddots & 0 \\ 0 & \dots & 0 & s_{k,2\delta-1} \end{pmatrix},$$

where $F_\alpha \in \mathbb{C}^{\alpha \times \alpha}$ is the unitary $\alpha \times \alpha$ discrete Fourier transform matrix, and $\{s_{k,j}\}_{k \in \{1,2,\dots,d\}}^{j \in \{1,2,\dots,2\delta-1\}}$ can be explicitly evaluated.

Since $F_{2\delta-1}$ is unitary,

$$\min_{j \in [2\delta-1]} |s_{k,j}| \leq \sigma_{2\delta-1}(J_k) \leq \sigma_1(J_k) \leq \max_{j \in [2\delta-1]} |s_{k,j}|.$$

Appendix: Condition Number Proof Sketch

(Step 3) Bound the maximum and minimum values of $|s_{k,j}|$ from above and below, respectively, over all $k \in \{1, 2, \dots, d\}$ and $j \in \{1, 2, \dots, 2\delta - 1\}$. Minimize upper bound w.r.t. a parameter.

(Step 4) Final result obtained using

$$\kappa(M') = \frac{\sigma_1(M')}{\sigma_D(M')} = \frac{\sigma_1(J)}{\sigma_D(J)} \leq \frac{\max_{k \in \{1, 2, \dots, d\}} \sigma_1(J_k)}{\min_{k \in \{1, 2, \dots, d\}} \sigma_{2\delta-1}(J_k)}.$$

Appendix: Flattening “Non-Sparse” Vectors

- Recall: Due to compact support of our masks, only $\lfloor \frac{\delta-3}{2} \rfloor$ -flat vectors can be recovered
- Arbitrary vectors can be “flattened” by multiplication with a random unitary matrix such as $W = PFB$, where
 - $P \in \{0, 1\}^{d \times d}$ is a permutation matrix selected uniformly at random from the set of all $d \times d$ permutation matrices
 - F is the unitary $d \times d$ discrete Fourier transform matrix
 - $B \in \{-1, 1\}^{d \times d}$ is a random diagonal matrix with i.i.d. symmetric Bernoulli entries on its diagonal

Some Details

Definition

Let $m \in \{1, 2, \dots, d\}$. A vector $\mathbf{x} \in \mathbb{C}^d$ will be called m -flat if its entries can be partitioned into at least $\lfloor \frac{d}{m} \rfloor$ contiguous blocks such that:

- 1 Every block contains either m or $m + 1$ entries,
- 2 Every block contains at least one entry whose magnitude is $\geq \frac{\|\mathbf{x}\|_2}{2\sqrt{d}}$, and

– The smaller m is, the flatter (and less sparse) \mathbf{x} must be.

Some Details

Definition

Let $\epsilon \in (0, 1)$, and $S \subset \mathbb{C}^d$ be finite. An $m \times d$ matrix A is a linear Johnson-Lindenstrauss embedding of S into \mathbb{C}^m if

$$(1 - \epsilon) \| \mathbf{u} - \mathbf{v} \|_2^2 \leq \| A\mathbf{u} - A\mathbf{v} \|_2^2 \leq (1 + \epsilon) \| \mathbf{u} - \mathbf{v} \|_2^2$$

holds $\forall \mathbf{u}, \mathbf{v} \in S \cup \{\mathbf{0}\}$. In this case we will say that A is a $JL(m, d, \epsilon)$ -embedding of S into \mathbb{C}^m .

JL embeddings are closely related to the *Restricted Isometry Property (RIP)*. A matrix with the restricted isometry property can be used to construct a Johnson-Lindenstrauss embedding matrix.

Some Details

- $W = PFB$
- For any given $m \in \{1, 2, \dots, d\}$, one can partition W into $\lfloor \frac{d}{m} \rfloor$ blocks of contiguous rows,

$$W = \left(W_1 W_2 \dots W_{\lfloor \frac{d}{m} \rfloor} \right)^T.$$

- Each renormalized sub-matrix of W , $\sqrt{\frac{d}{m}} \cdot W_j$ is “almost” a random sampling matrix times a random diagonal Bernoulli matrix and behaves like a $\text{JL}(m, d, \epsilon)$ -embedding of our signal \mathbf{x} into \mathbb{C}^m (or \mathbb{C}^{m+1}).
- Each block of m consecutive entries of $W\mathbf{x}$ should have roughly the same ℓ_2 -norm as one another.